

Convex Optimization and Duality

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Based on the Online Course "Convex Optimization" by Stephen Boyd and the Book "Convex Optimization" by Stephen Boyd and Lieven Vandenberghe

Sources

Book:

http://web.stanford.edu/~boyd/cvxbook/bv_cvxbook.pdf
MOOC:

https://lagunita.stanford.edu/courses/Engineering/CVX101/Winter2014/about

- ► Theory: 8h40 + exercises.
 - Convex sets,
 - Convex functions,
 - Convex optimization problems,
 - Duality.
- ► Applications: 3h15 + exercises.
 - Approximation and fitting, Statistical estimation, Geometric problems.
- ► Algorithms: 5h15 + exercises.
 - Numerical linear algebra, Unconstrained minimization, Equality constrained minimization, Interior point methods.



Introduction

Plan

Convex Sets

Convex Functions

Optimization

Optimality Conditions



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Fundamental Idea

A convex set \mathcal{C} can be fully described by its supporting hyperplanes.





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$$\begin{split} \phi_{\boldsymbol{\lambda}} &= \inf\{\boldsymbol{\lambda}^{T}\boldsymbol{x}, \boldsymbol{x} \in \mathcal{C}\} \\ &= \sup\{\phi \mid \forall \boldsymbol{x} \in \mathcal{C}, \boldsymbol{\lambda}^{T}\boldsymbol{x} \geq \phi\}. \end{split}$$





Fundamental Idea

1

A convex set \mathcal{C} can be fully described by its supporting hyperplanes.

$$\begin{split} \phi_{\boldsymbol{\lambda}} &= \inf\{\boldsymbol{\lambda}^{T} \boldsymbol{x}, \boldsymbol{x} \in \mathcal{C}\}\\ &= \sup\{\phi \mid \forall \boldsymbol{x} \in \mathcal{C}, \boldsymbol{\lambda}^{T} \boldsymbol{x} \geq \phi\}. \end{split}$$

What is the dual space? Example:

- ▶ x₁, x₂... in (say) kg,
- ▶ φ in (say) \$,
- ► Then $\lambda_1, \lambda_2, \dots$ in $\$.kg^{-1}$.

 $\Rightarrow \text{ It is a space of unit prices.} \\\Rightarrow \text{ If the unit prices were } \boldsymbol{\lambda} \text{, then each bundle in } \mathcal{C} \text{ would cost at least } \phi_{\boldsymbol{\lambda}} \text{.}$





Separating Hyperplanes Theorem

If two convex sets do not intersect, then they can be separated by (at least) a hyperplane.







Convex Sets

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Convex Functions $\mathbb{R}^n \to \mathbb{R}$

Definition:

- And **dom** *f* is a convex set.
- I.e. the *epigraph* of f is a convex set.



Concave: -f is convex. I.e. the *hypograph* of f is a convex set.

Affine \Leftrightarrow convex and concave.



Convex Functions $\mathbb{R}^n \to \mathbb{R}$: Differentiable Case

If differentiable:

$$f(\mathbf{y}) \geq f(\mathbf{x}) + \nabla f(\mathbf{x})^T (\mathbf{y} - \mathbf{x}).$$

 \Rightarrow First-order Taylor gives a global underestimator.

$$\Rightarrow$$
 For example, if $oldsymbol{
abla} f(oldsymbol{x}) = oldsymbol{0}$. .



If differentiable twice: $\nabla^2 f(\mathbf{x}) \succeq 0$ (the Hessian is semidefinite positive).



Examples of Convex Functions

- ► Any norm on \mathbb{R}^n ,
- ▶ $\mathbf{x} \rightarrow \mathbf{x}^T P \mathbf{x}$, where P is semidefinite positive,
- $\blacktriangleright \ \boldsymbol{x} \to \max_i(x_i),$
- ▶ $\mathbf{x} \rightarrow \log \sum_{i} \exp(x_i)$ ("LogSumExp"),
- $\mathbf{x} \rightarrow \left(\prod_{i} x_{i}\right)^{\frac{1}{n}}$ (geometric mean),
- ▶ $\textbf{x} \to \infty_{\textbf{x} \notin \mathcal{C}}$, where \mathcal{C} is a convex set,
- ▶ $X \rightarrow \log \det X$ over the set of definite positive matrices,
- $X \rightarrow \text{eigenvalue}_{\max}(X)$ over the set of symmetric matrices.



Convex Functions Calculus

Are convex:

- $\blacktriangleright \mathbf{x} \to f(\mathbf{x}) + g(\mathbf{x}),$
- $\blacktriangleright \ \textbf{\textit{x}} \to \lambda f(\textbf{\textit{x}}), \ \text{for} \ \lambda \geq 0,$
- $\blacktriangleright x \to f(Ax + b),$
- $\blacktriangleright \ \mathbf{x} \to sup_{\theta \in \Theta} f_{\theta}(\mathbf{x}),$



Composition rule. Assume:

- Functions v_i are convex, c_i concave, a_i affine,
- f convex, nondecreasing in each argument v_i , nonincreasing in each argument c_i . Then $\mathbf{x} \to f(v_1(\mathbf{x}), \ldots, c_1(\mathbf{x}), \ldots, a_1(\mathbf{x}), \ldots)$ is convex.

[Mnemonic: write second derivative. But still true if not differentiable.]



Prove that a Function is Convex in Practice

- Apply the definition (extremely rare),
- Compute the Hessian / second derivative (to be avoided if possible),
- ▶ Prove that any restriction to a line, i.e. $t \rightarrow f(\mathbf{x} + t\mathbf{v})$, is convex (sometimes),
- Apply the rules of the previous slide (laziest hence preferred method).



Sublevel Sets

If *f* is convex, then:

$$\mathcal{C}_{\alpha} = \{ \boldsymbol{x} \in \operatorname{\mathsf{dom}} f \mid f(\boldsymbol{x}) \leq \alpha \}$$

is convex.

 \Rightarrow Often used to prove that a set is convex.





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Optimization Problem in the Standard Form

$$\begin{array}{ll} \text{minimize} & f(\boldsymbol{x}) \\ \text{subject to} & g_i(\boldsymbol{x}) \leq 0, \quad i = 1, \dots, r, \\ & h_j(\boldsymbol{x}) = 0, \quad j = 1, \dots, s. \end{array}$$

•
$$f(x)$$
: cost (e.g. in \$).

Example of an inequality constraint:

$$v_1x_1+v_2x_2-S\leq 0$$

where x_i in kg, v_i in m³.kg⁻¹ and S is the volume of my warehouse in m³.

 (Theoretical) remark: any equality constraint can be seen as two opposite inequality constraints.



$\begin{array}{ll} \mbox{minimize} & f(\textbf{\textit{x}}) \\ \mbox{subject to} & g_i(\textbf{\textit{x}}) \leq 0, & i = 1, \ldots, r, \\ & h_j(\textbf{\textit{x}}) = 0, & j = 1, \ldots, s. \end{array}$

A Bit of Vocabulary

$$\mathcal{D} = \operatorname{dom} f \cap \bigcap_{i} \operatorname{dom} g_{i} \cap \bigcap_{j} \operatorname{dom} h_{j}: \operatorname{domain} of the problem.$$

$$\mathcal{F} = \{ \mathbf{x} \in \mathcal{D} \mid \forall i, g_i(\mathbf{x}) \leq 0 \text{ and } \forall j, h_j(\mathbf{x}) = 0 \}$$
: set of *feasible points*.

$$p^* = \inf_{\mathbf{x} \in \mathcal{F}} f(\mathbf{x})$$
: optimal value of the problem.

▶
$$p^* = \infty$$
: the problem is *infeasible*, i.e. $\mathcal{F} = \emptyset$.

▶
$$p^* = -\infty$$
: the problem is *unbounded below*.

▶ *p*^{*} is finite:

- If $f(\mathbf{x}^*) = p^*$, then \mathbf{x}^* is an optimal point or solution.
- If no such x^* , then the optimal value p^* is not attained.



minimize	$f(\mathbf{x})$	
subject to	$g_i(\mathbf{x}) \leq 0,$	$i=1,\ldots,r,$
	$h_j(\mathbf{x}) = 0,$	$j=1,\ldots,s.$

Convex Optimization Problem

We say that the problem is convex if:

- f is convex,
- All g_i are convex,
- ▶ And all h_j are affine.

Motivation: this ensures that f restricted to \mathcal{F} is a convex function.

Then:

- ► Any local minimum is a global optimum, i.e. a solution.
- (Generally) efficient algorithms to find a solution.



minimize	$f(\mathbf{x})$	
subject to	$g_i(\mathbf{x}) \leq 0$,	$i=1,\ldots,r,$
	$h_j(\mathbf{x}) = 0,$	$j=1,\ldots,s.$

Lagrangian: Motivation We do not assume that the problem is convex.

The problem is equivalent to:

minimize
$$f(\mathbf{x}) + \sum_{i} \infty_{g_i(\mathbf{x}) > 0} + \sum_{j} \infty_{h_j(\mathbf{x}) \neq 0}$$
 $(\mathbf{x} \in D)$

Example with constraint $v_1x_1 + v_2x_2 - S \le 0$: using more than S has infinite cost, using less than S is costless (but does not generate an income).

Relaxation of the problem: fix a unit price $\lambda \ge 0$ (in $.m^{-3}$). Use more than S: buy space at price λ . Use less: sell the extra space at price λ .

minimize
$$f(\mathbf{x}) + \sum_{i} \lambda_{i} g_{i}(\mathbf{x}) + \sum_{j} \nu_{j} h_{j}(\mathbf{x})$$
 $(\mathbf{x} \in D)$



Optimization

minimize	$f(\mathbf{x})$	
subject to	$g_i(\mathbf{x}) \leq 0,$	$i=1,\ldots,r,$
	$h_j(\mathbf{x}) = 0,$	$j = 1, \ldots, s.$

Lagrangian and Dual Lagrangian

We do **not** assume that the problem is convex.

Lagrangian:

$$L(\boldsymbol{x},\boldsymbol{\lambda},\boldsymbol{\nu})=f(\boldsymbol{x})+\sum_{j}\lambda_{j}g_{j}(\boldsymbol{x})+\sum_{j}\nu_{j}h_{j}(\boldsymbol{x})$$

Dual Lagrangian:

$$\phi(\boldsymbol{\lambda},\boldsymbol{\nu}) = \inf_{\boldsymbol{x}\in\mathcal{D}} L(\boldsymbol{x},\boldsymbol{\lambda},\boldsymbol{\nu})$$

It is always concave (even if the original problem is not convex).

▶ It provides a *parametrized family of lower bounds* for *f*. Indeed, for any feasible *x* and any $\lambda \ge 0$ (and without any requirement on ν):

$$\phi(\boldsymbol{\lambda},\boldsymbol{\nu}) \leq L(\boldsymbol{x},\boldsymbol{\lambda},\boldsymbol{\nu}) = f(\boldsymbol{x}) + \sum_{i} \underbrace{\lambda_{i}g_{i}(\boldsymbol{x})}_{\leq 0} + \sum_{j} \nu_{j} \underbrace{h_{j}(\boldsymbol{x})}_{=0} \leq f(\boldsymbol{x}).$$



Dual Problem

We do **not** assume that the problem is convex.

We have $\phi(\lambda, \nu) \leq p^*$. To find the best lower bound, let us maximize ϕ . This is the dual problem:

maximize
$$\phi(\boldsymbol{\lambda}, \boldsymbol{\nu})$$

subject to $\lambda_i \geq 0$, $i = 1, \dots, r$.

This is a convex problem, hence (generally) convenient to solve, at least numerically.

Its optimal value is denoted by d^* . By construction, we have $d^* \le p^*$. Duality gap: $p^* - d^*$.

- If = 0: "strong duality" situation.
- If > 0: "weak duality" situation.



Slater's Conditions

 $\begin{array}{ll} \mbox{minimize} & f(\mathbf{x}) \\ \mbox{subject to} & g_i(\mathbf{x}) \leq 0, & i = 1, \dots, r, \\ & h_j(\mathbf{x}) = 0, & j = 1, \dots, s. \end{array} \\ L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f(\mathbf{x}) + \sum_i \lambda_i g_i(\mathbf{x}) + \sum_j \nu_j h_j(\mathbf{x}) \\ \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_x L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{maximize} & \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{subject to} & \lambda_i \geq 0, & i = 1, \dots, r. \end{array}$

lf:

- The primal problem is convex,
- And the constraints are "strictly feasible" (i.e. with strict inequalities), Then:

►
$$p^{\star} = d^{\star}$$
,

• If $p^* = d^* > -\infty$, then the dual optimum is attained, i.e. $\exists (\lambda^*, \nu^*)$ s.t. $\phi(\lambda^*, \nu^*) = d^* = p^*$.

Remark: for *affine* inequality constraint, the "strictness" condition can be dropped.



Geometric Interpretation: Set of Values

We do **not** assume that the problem is convex.

Example: only one constraint $g(\mathbf{x}) \leq 0$.

$$egin{aligned} \mathcal{G} &= \Big\{ ig(oldsymbol{x}), f(oldsymbol{x})ig) \mid oldsymbol{x} \in \mathcal{D} \Big\}. \ oldsymbol{p}^{\star} &= \inf\{t \mid (u,t) \in \mathcal{G} ext{ and } u \leq 0\}. \end{aligned}$$

$$egin{aligned} \phi(\lambda) &= \inf\{f(m{x}) + \lambda g(m{x}) \mid m{x} \in \mathcal{D}\} \ &= \inf\{\lambda u + t \mid (u, t) \in \mathcal{G}\} \ &= \sup\{b \mid orall (u, t) \in \mathcal{G}, \lambda u + t \geq b\} \end{aligned}$$

 $\Rightarrow \lambda u + t \ge \phi(\lambda) \text{ is a supporting}$ hyperplane of \mathcal{G} (of slope $-\lambda$).



$$\begin{array}{lll} \mbox{minimize} & f(\mathbf{x}) \\ \mbox{subject to} & g_i(\mathbf{x}) \leq 0, & i = 1, \dots, r, \\ & h_j(\mathbf{x}) = 0, & j = 1, \dots, s. \end{array} \\ L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f(\mathbf{x}) + \sum_i \lambda_i g_i(\mathbf{x}) + \sum_j \nu_j h_j(\mathbf{x}) \\ \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_x L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{maximize} & \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{subject to} & \lambda_i \geq 0, & i = 1, \dots, r. \end{array}$$



Geometric Interpretation: Epigraph variation

We do **not** assume that the problem is convex.

Example: only one constraint $g(\mathbf{x}) \leq 0$.

- $\begin{aligned} \mathcal{A} &= \big\{ (u,t) \mid \exists \pmb{x} \in \mathcal{D}, u \geq \pmb{g}(\pmb{x}), t \geq f(\pmb{x}) \big\} \\ &= \mathcal{G} \ \cup \text{ points that are worse.} \end{aligned}$
- $p^{\star} = \inf\{t \mid (0, t) \in \mathcal{A}\}.$

$$egin{aligned} \phi(\lambda) &= \inf\{f(m{x}) + \lambda g(m{x}) \mid m{x} \in \mathcal{D}\} \ &= \inf\{\lambda u + t \mid (u, t) \in \mathcal{A}\} \ &= \sup\{b \mid orall (u, t) \in \mathcal{A}, \lambda u + t \geq b\}. \end{aligned}$$

 $\Rightarrow \lambda u + t \ge \phi(\lambda) \text{ is a supporting}$ hyperplane of \mathcal{A} (of slope $-\lambda$).



$$\begin{array}{ll} \begin{array}{ll} \mbox{minimize} & f(\mathbf{x}) \\ \mbox{subject to} & g_i(\mathbf{x}) \leq 0, & i=1,\ldots,r, \\ & h_j(\mathbf{x}) = 0, & j=1,\ldots,s. \end{array} \\ L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f(\mathbf{x}) + \sum_i \lambda_i g_i(\mathbf{x}) + \sum_j \nu_j h_j(\mathbf{x}) \\ \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) = inf_{\boldsymbol{\lambda}} L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{maximize} & \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{subject to} & \lambda_i \geq 0, \quad i=1,\ldots,r. \end{array}$$



Convex Problems: Why There is (Usually) Strong Duality

If the optimization problem is convex, then $\ensuremath{\mathcal{A}}$ is convex.

There is a separating hyperplane between \mathcal{A} and $\mathcal{B} = \{(0, t) \mid t < p^{\star}\}.$

[Here, constraints qualification such as Slater's conditions ensure that the hyperplane is not vertical.]

This hyperplane gives the good λ and $\phi(\lambda)$.

 $\begin{array}{ll} \mbox{minimize} & f(\mathbf{x}) \\ \mbox{subject to} & g_i(\mathbf{x}) \leq 0, & i = 1, \dots, r, \\ & h_j(\mathbf{x}) = 0, & j = 1, \dots, s. \\ \mbox{} L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f(\mathbf{x}) + \sum_i \lambda_i g_i(\mathbf{x}) + \sum_j \nu_j h_j(\mathbf{x}) \\ \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_x L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{maximize} & \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{subject to} & \lambda_i \geq 0, & i = 1, \dots, r. \end{array}$





Example of Convex Problem Without Strong Duality

 $\begin{array}{ll} \text{minimize} & e^{-x} \\ \text{subject to} & x^2/y \leq 0 \\ \text{with } \mathcal{D} = \{(x,y) \mid y > 0\}. \end{array}$

This is a convex problem. Constraint means x = 0. N.B.: Slater's conditions are violated.

$$p^{\star} = 1$$

$$\phi(\lambda) = \inf_{x,y} e^{-x} + \lambda \frac{x^2}{y} = 0$$

$$d^{\star} = 0$$

$$\Rightarrow \text{ No strong duality.}$$

 $\begin{array}{lll} & \mbox{minimize} & f(\mathbf{x}) \\ & \mbox{subject to} & g_i(\mathbf{x}) \leq 0, & i = 1, \dots, r, \\ & h_j(\mathbf{x}) = 0, & j = 1, \dots, s. \\ & L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f(\mathbf{x}) + \sum_i \lambda_i g_i(\mathbf{x}) + \sum_j \nu_j h_j(\mathbf{x}) \\ & \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_{\boldsymbol{\lambda}} L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \\ & \mbox{maximize} & \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) \\ & \mbox{subject to} & \lambda_i \geq 0, & i = 1, \dots, r. \end{array}$





Saddle-point interpretation

We do **not** assume that the problem is convex.

$$\begin{split} \sup_{\boldsymbol{\lambda} \ge 0, \boldsymbol{\nu}} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) &= \sup_{\boldsymbol{\lambda} \ge 0, \boldsymbol{\nu}} \left(f(\boldsymbol{x}) + \sum_{i} \lambda_{i} g_{i}(\boldsymbol{x}) + \sum_{j} \nu_{j} h_{j}(\boldsymbol{x}) \right) \\ &= \left| \begin{array}{c} f(\boldsymbol{x}) \text{ if } \boldsymbol{x} \in \mathcal{F}, \\ \infty \text{ otherwise.} \end{array} \right. \end{split}$$

$$\begin{array}{ll} \begin{array}{ll} \mbox{minimize} & f(\mathbf{x}) \\ \mbox{subject to} & g_i(\mathbf{x}) \leq 0, & i=1,\ldots,r, \\ & h_j(\mathbf{x}) = 0, & j=1,\ldots,s. \end{array} \\ L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f(\mathbf{x}) + \sum_i \lambda_i g_i(\mathbf{x}) + \sum_j \nu_j h_j(\mathbf{x}) \\ \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) = inf_x L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{maximize} & \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{subject to} & \lambda_i \geq 0, \quad i=1,\ldots,r. \end{array}$$

$$\Rightarrow \rho^{\star} = \inf_{\boldsymbol{x}\in\mathcal{D}} \sup_{\boldsymbol{\lambda}\geq 0, \boldsymbol{
u}} L(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{
u}).$$

And by definition: $d^{\star} = \sup_{\lambda \geq 0, \nu} \inf_{\mathbf{x} \in \mathcal{D}} L(\mathbf{x}, \lambda, \nu).$

Hence weak duality is simply a particular case of the max-min inequality: $\sup_{y} \inf_{x} f(x, y) \leq \inf_{x} \sup_{y} f(x, y)$. And:

 $(\textbf{\textit{x}},(\boldsymbol{\lambda},\boldsymbol{\nu})) \text{ saddle-point of } L \Leftrightarrow \textbf{\textit{x}} = \textbf{\textit{x}}^{\star}, \boldsymbol{\lambda} = \boldsymbol{\lambda}^{\star}, \boldsymbol{\nu} = \boldsymbol{\nu}^{\star} \text{ and } L(\textbf{\textit{x}},\boldsymbol{\lambda},\boldsymbol{\nu}) = p^{\star} = d^{\star}.$



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Complementary slackness

We do **not** assume that the problem is convex.

 $\begin{array}{ll} \mbox{minimize} & f(\mathbf{x}) \\ \mbox{subject to} & g_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, r, \\ & h_j(\mathbf{x}) = 0, \quad j = 1, \dots, s. \end{array} \\ L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f(\mathbf{x}) + \sum_i \lambda_i g_i(\mathbf{x}) + \sum_j \nu_j h_j(\mathbf{x}) \\ \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_x L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{maximize} & \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{subject to} & \lambda_i \geq 0, \quad i = 1, \dots, r. \end{array}$

Assume strong duality: $d^* = p^*$. Assume these optimal values are reached at (λ^*, ν^*) and \mathbf{x}^* respectively. Then all the following inequalities are in fact equalities:

$$\phi(\boldsymbol{\lambda}^{\star}, \boldsymbol{\nu}^{\star}) \leq L(\boldsymbol{x}^{\star}, \boldsymbol{\lambda}^{\star}, \boldsymbol{\nu}^{\star}) = f(\boldsymbol{x}^{\star}) + \sum_{i} \underbrace{\lambda_{i}^{\star} g_{i}(\boldsymbol{x}^{\star})}_{\leq 0} + \sum_{j} \nu_{j}^{\star} \underbrace{h_{j}(\boldsymbol{x}^{\star})}_{=0} \leq f(\boldsymbol{x}^{\star}).$$

Hence:

►
$$L(\mathbf{x}^{\star}, \mathbf{\lambda}^{\star}, \mathbf{\nu}^{\star}) = f(\mathbf{x}^{\star}) = \phi(\mathbf{\lambda}^{\star}, \mathbf{\nu}^{\star}),$$

► $\forall i, \lambda_i^{\star} g_i(\mathbf{x}^{\star}) = 0$ (complementary slackness). I.e. $\lambda_i^{\star} = 0$ or $g_i(\mathbf{x}^{\star}) = 0.$



Karush-Kuhn Tucker conditions

We do **not** assume that the problem is convex.

We now assume that f, all g_i and all h_j are differentiable.

1. If
$$f(\mathbf{x}^*) = p^* = d^* = \phi(\lambda^*, \nu^*)$$
, then (KKT conditions):

►
$$\forall i, g_i(\mathbf{x}^*) \leq 0$$
 and $\forall j, h_j(\mathbf{x}^*) = 0$ (primal feasibility),

▶
$$\forall i, \lambda_i^{\star} \geq 0$$
 (dual feasibility),

►
$$\forall i, \lambda_i^* g_i(\mathbf{x}^*) = 0$$
 (complementary slackness).

2. If the problem is convex, then the converse is true.

3. If the problem is convex and satisfies Slater's conditions, then x^* is optimal iff there exists (λ^*, ν^*) that meets KKT conditions.



$$\begin{array}{lll} & \text{minimize} & f(\mathbf{x}) \\ & \text{subject to} & g_i(\mathbf{x}) \leq 0, & i = 1, \dots, r, \\ & h_j(\mathbf{x}) = 0, & j = 1, \dots, s. \end{array} \\ & \mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f(\mathbf{x}) + \sum_i \lambda_i g_i(\mathbf{x}) + \sum_j \nu_j h_j(\mathbf{x}) \\ & \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_{\boldsymbol{x}} \mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \\ & \text{maximize} & \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) \\ & \text{subject to} & \lambda_i \geq 0, & i = 1, \dots, r. \end{array}$$

KKT: Geometric Interpretation (One Constraint)

minimize
$$f(\mathbf{x})$$
Example:subject to $g_1(\mathbf{x}) \leq 0$ $x_1 - x_2 \leq 0$

If x^* is in the interior:

 $\boldsymbol{\nabla} f(\boldsymbol{x}^{\star}) = \boldsymbol{0}$





KKT: Geometric Interpretation (One Constraint)

$$\begin{array}{lll} \mbox{minimize} & f({\pmb x}) & \mbox{Example:} \\ \mbox{subject to} & g_1({\pmb x}) \leq 0 & & x_1 - & x_2 \leq 0 \end{array}$$

If x^* is in the interior:

$$\boldsymbol{\nabla} f(\boldsymbol{x}^{\star}) = \boldsymbol{0}$$

If
$$x^*$$
 is on the frontier:

$$- oldsymbol{
abla} f(oldsymbol{x}^{\star}) = \lambda_1 oldsymbol{
abla} g_1(oldsymbol{x}^{\star}) \quad (ext{with } \lambda_1 \geq 0)$$





KKT: Geometric Interpretation (One Constraint)

$$\begin{array}{lll} \mbox{minimize} & f({\pmb x}) & \mbox{Example:} \\ \mbox{subject to} & g_1({\pmb x}) \leq 0 & & x_1 - & x_2 \leq 0 \end{array}$$

If \mathbf{x}^{\star} is in the interior:

$$\boldsymbol{\nabla} f(\boldsymbol{x}^{\star}) = \boldsymbol{0}$$

If \mathbf{x}^{\star} is on the frontier:

$$- oldsymbol{
abla} f(oldsymbol{x}^{\star}) = \lambda_1 oldsymbol{
abla} g_1(oldsymbol{x}^{\star}) \quad (ext{with } \lambda_1 \geq 0)$$

Anyway, the second condition is met. Moreover, if $\lambda_1 > 0$, then $g_1(\mathbf{x}) = 0$.





KKT: Geometric Interpretation (Several Constraints)

 $\begin{array}{ll} \text{minimize} & f({\pmb{x}}) \\ \text{subject to} & g_1({\pmb{x}}) \leq 0 \\ & g_2({\pmb{x}}) \leq 0 \end{array}$

 \mathbf{x}^{\star} interior: $\nabla f(\mathbf{x}^{\star}) = \mathbf{0}$. \mathbf{x}^{\star} on first frontier: $-\nabla f(\mathbf{x}^{\star}) = \lambda_1 \nabla g_1(\mathbf{x}^{\star})$. \mathbf{x}^{\star} on second frontier: $-\nabla f(\mathbf{x}^{\star}) = \lambda_2 \nabla g_2(\mathbf{x}^{\star})$.





KKT: Geometric Interpretation (Several Constraints)

 $\begin{array}{ll} \text{minimize} & f(\boldsymbol{x}) \\ \text{subject to} & g_1(\boldsymbol{x}) \leq 0 \\ & g_2(\boldsymbol{x}) \leq 0 \end{array}$

 \mathbf{x}^{\star} interior: $\nabla f(\mathbf{x}^{\star}) = \mathbf{0}$. \mathbf{x}^{\star} on first frontier: $-\nabla f(\mathbf{x}^{\star}) = \lambda_1 \nabla g_1(\mathbf{x}^{\star})$. \mathbf{x}^{\star} on second frontier: $-\nabla f(\mathbf{x}^{\star}) = \lambda_2 \nabla g_2(\mathbf{x}^{\star})$. If \mathbf{x}^{\star} is on the intersection of frontiers:

$$- \boldsymbol{\nabla} f(\boldsymbol{x}^{\star}) = \lambda_1 \boldsymbol{\nabla} g_1(\boldsymbol{x}^{\star}) + \lambda_2 \boldsymbol{\nabla} g_2(\boldsymbol{x}^{\star}) \quad (\text{with } \lambda_i \geq 0).$$

In all cases, the above condition is met. Moreover, if $\lambda_i > 0$, then $g_i(\mathbf{x}) = 0$. Example: $x_1 - x_2 < 0$ $-2x_1 + 0.25x_2 < 0$ X2 ∇g_2 χ́1



Sensitivity Analysis (Very Quickly)

$$\begin{array}{ll} \text{minimize} & f(\boldsymbol{x}) \\ \text{subject to} & g_i(\boldsymbol{x}) \leq u_i, \quad i = 1, \dots, r, \\ & h_j(\boldsymbol{x}) = v_j, \quad j = 1, \dots, s. \end{array}$$

Denote $p^*(\boldsymbol{u}, \boldsymbol{v})$ its optimal value.

If strong duality holds and if p^* is differentiable at (0, 0), then:

$$\lambda_i^{\star} = -\frac{\partial p^{\star}}{\partial u_i}(\mathbf{0},\mathbf{0}) \text{ and } \nu_i^{\star} = -\frac{\partial p^{\star}}{\partial v_i}(\mathbf{0},\mathbf{0}).$$



$$\begin{array}{lll} \mbox{minimize} & f(\mathbf{x}) \\ \mbox{subject to} & g_i(\mathbf{x}) \leq 0, & i = 1, \dots, r, \\ & h_j(\mathbf{x}) = 0, & j = 1, \dots, s. \end{array} \\ \mathcal{L}(\mathbf{x}, \mathbf{\lambda}, \boldsymbol{\nu}) = f(\mathbf{x}) + \sum_i \lambda_i g_i(\mathbf{x}) + \sum_j \nu_j h_j(\mathbf{x}) \\ & \phi(\mathbf{\lambda}, \boldsymbol{\nu}) = inf_x \mathcal{L}(\mathbf{x}, \mathbf{\lambda}, \boldsymbol{\nu}) \\ & \mbox{maximize} & \phi(\mathbf{\lambda}, \boldsymbol{\nu}) \\ & \mbox{subject to} & \lambda_i \geq 0, & i = 1, \dots, r. \end{array}$$

Take-aways

- $\begin{array}{lll} \mbox{minimize} & f(\mathbf{x}) \\ \mbox{subject to} & g_i(\mathbf{x}) \leq 0, & i = 1, \dots, r, \\ & h_j(\mathbf{x}) = 0, & j = 1, \dots, s. \end{array} \\ L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f(\mathbf{x}) + \sum_i \lambda_i g_i(\mathbf{x}) + \sum_j \nu_j h_j(\mathbf{x}) \\ \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_x L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{maximize} & \phi(\boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \mbox{subject to} & \lambda_i \geq 0, & i = 1, \dots, r. \end{array}$
- The Lagrangian describes a *relaxation* of the problem with a *unit price* for each constraint (Lagrange multiplier).
- The dual Lagrangian provides a parametrized family of lower bounds for the primal problem.
- ▶ The **dual problem** is *always convex*, even if the primal problem is not.
- When the primal problem is convex, there is usually strong duality (with mild additional assumptions such as *Slater's conditions*).
- For differentiable problems, think of KKT conditions (*necessary* if there is strong duality, *sufficient* if the problem is convex).





Questions?

